

EE 2372 Test 2

9 problems, 100 points.

November 24, 1998

SOLUTIONS

NAME

Closed book, closed notes, no calculators. Scratch paper will be provided, so do not use any of your own.

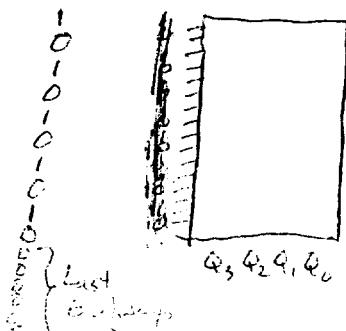
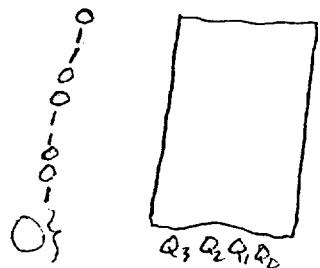
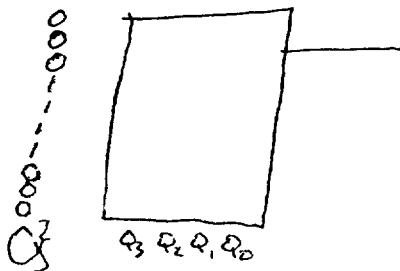
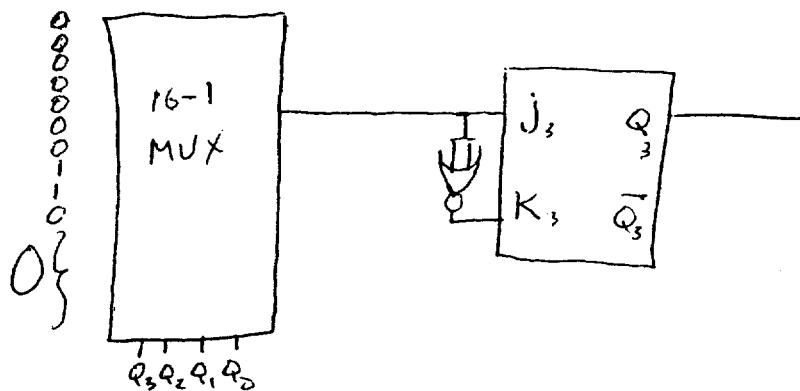
You are permitted pens or pencils, erasers, and a (non-calculator) watch. All other items are to be placed underneath your desk.

Please read the entire exam before beginning, and note point values. Some problems are more worthwhile than others.

Do not turn this page until instructed to do so.

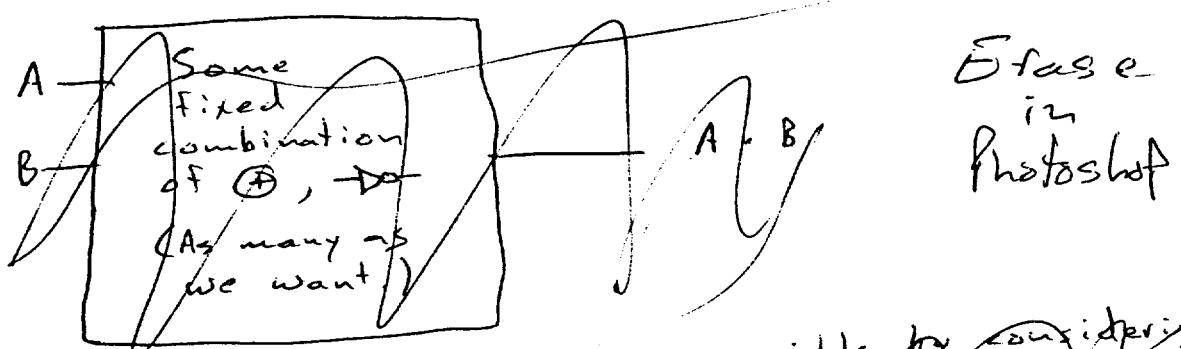
Good luck!

1. Design a synchronous decade (0-9) counter using four JK flip-flops, four NOR gates, and four 16-1 multiplexers. Any erroneous (glitch) state must be reset to zero on the next clock cycle. The states are $(Q_3 Q_2 Q_1 Q_0)$
 $\text{(MSB) } \text{(LSB)}$



2. No, we can't make any gate we want, using only XOR and inverters.

If we could we could make an AND gate. It would look like this:



We can see this is impossible by considering the case of $A \text{ AND } A$. Of course $A \oplus A = 0$, but consider what can be done with XOR and inverters.

We have these possibilities

$$\begin{array}{c} A \\ B \end{array} \rightarrow \text{XOR gate} \rightarrow \overline{A \oplus B} = AB + \overline{A}\overline{B}$$

$$\begin{array}{c} A \\ B \end{array} \rightarrow \text{XOR gate} \rightarrow A \oplus \overline{B} = AB + \overline{A}\overline{B}$$

$$\begin{array}{c} A \\ B \end{array} \rightarrow \text{XOR gate} \rightarrow \overline{A \oplus \overline{B}} = \overline{AB} + A\overline{B}$$

$$\begin{array}{c} A \\ B \end{array} \rightarrow \text{XOR gate} \rightarrow \overline{A \oplus B} = AB + \overline{A}\overline{B}$$

These all reduce to $\overline{\oplus}$ or \oplus and we're stuck in a closed set of possibilities.

3.

J, K	Q^{t+1}	Q^t
00	0	0
01	0 (reset)	1
10	1 (set)	0
11	toggle	1

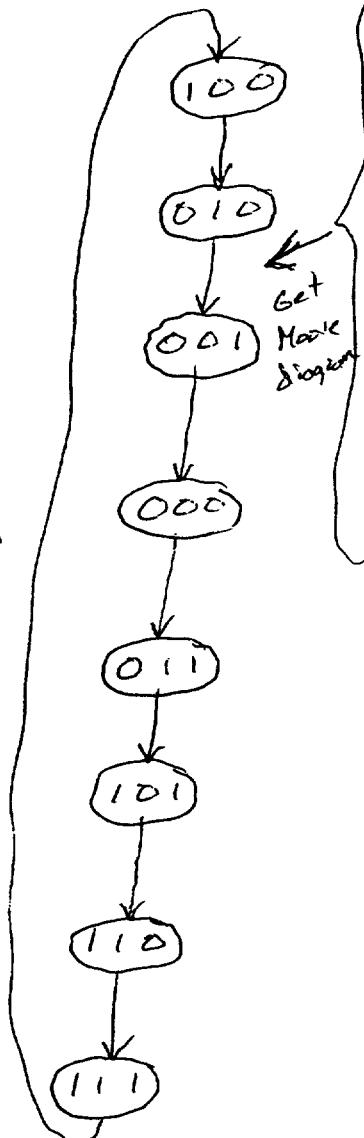
Q^t	Q^{t+1}	J, K
00	0	0X
01	1	1X
10	0	X1
11	0	X0

solve backwards
to get
this

cur		next		J_2, K_2	J, K_1	No. to
$Q_2 Q_1 Q_0$	$Q_2 Q_1 Q_0$	$Q_2 Q_1 Q_0$	$Q_2 Q_1 Q_0$	J_2, K_2	J, K_1	$Q_2 Q_1 Q_0$
000	011	011	010	0X	1X	
001	000	000	000	0X	0X	
010	001	001	0X	X1		
011	101	101	1X	X1		
100	010	010	X1	1X		
101	110	110	X0	1X		
110	111	111	X0	X0		
111	100	100	X0	X1		

Like
a
twisted
ring
counter
with
transition
states

(000, 111)



$$J_2 = Q_1 Q_0$$

Q_1	00	01	11	10
0	0	0	1	0
1	X	X	(X)	X

$$K_2 \neq \bar{Q}_1 \bar{Q}_0$$

X	X	X	X
1	0	0	0

$$J_1 = \bar{Q}_0 + Q_2$$

1	0	X	X
(1)	1	X	(X)

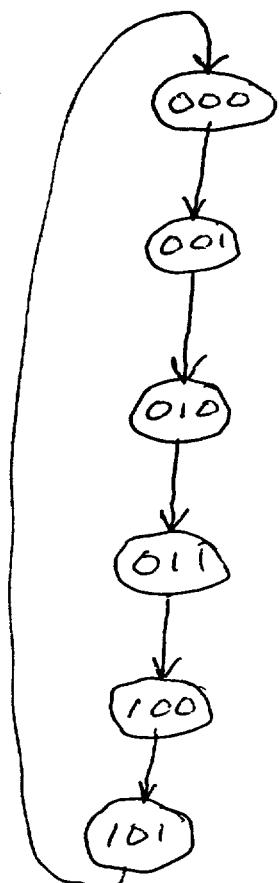
$$K_1 = \bar{Q}_2 + Q_0$$

X	X	1	0
X	X	1	0

$$J_0 = \bar{Q}_2 \bar{Q}_0 + \bar{Q}_2 Q_1 + Q_0 \bar{Q}_1$$

Q_2	00	01	11	10
0	D	O	1	1
1	O	O	0	0

4. Design a 3-bit synchronous Modulo-6 counter, using D flip flops: $D_2 D_1 D_0$.
 (It counts 0-5 and starts over.)
 Use a Moore diagram



$Q_2 Q_1 Q_0$	$D_2 D_1 D_0$
000	001
001	010
010	011
011	100
100	101
101	000
110	XX
111	XXX

$$D_0 = Q_1 Q_0 + Q_2 \bar{Q}_0$$

D_2	Q_2	$Q_1 Q_0$	00	01	11	10
0	0	0	0	0	1	1
1	1	1	0	0	X	X

$$D_1 = \bar{Q}_2 \bar{Q}_1 Q_0 + Q_1 \bar{Q}_0$$

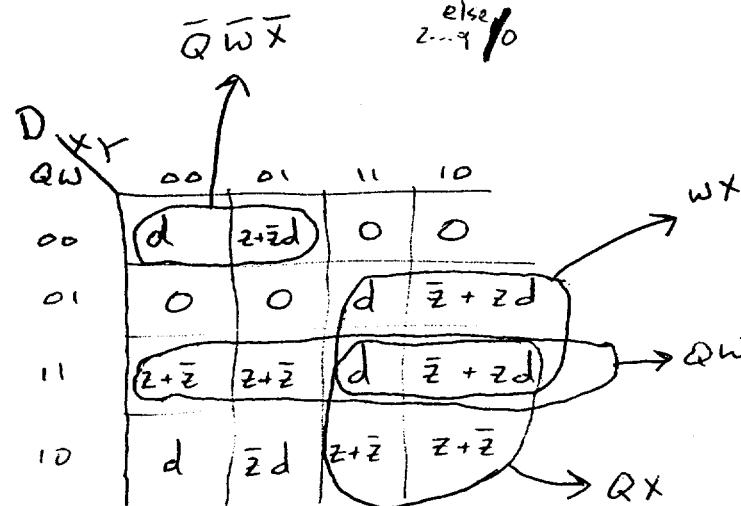
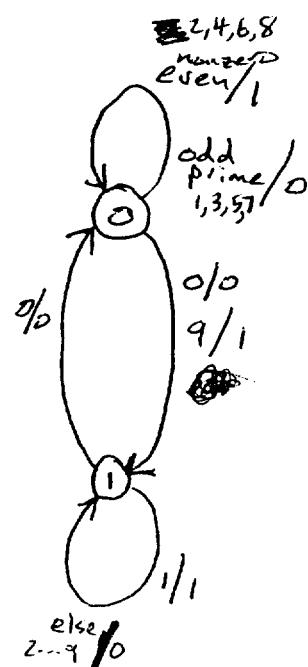
D_2	Q_2	$Q_1 Q_0$	00	01	11	10
0	0	0	0	0	0	1
1	1	1	0	0	X	X

$$D_0 = \bar{Q}_0$$

D_2	Q_2	$Q_1 Q_0$	00	01	11	10
0	1	0	1	0	0	1
1	0	1	0	0	X	X

$W'X'Y'Z'$	D	S	$D(z)$	$S(z)$
00000	x	x	d	$\bar{z}d$
00001	x	x		
00100	x	x	$z+\bar{z}d$	$\bar{z}d$
00111	1	0		
01000	0	0		
01001	0	1	0	z
01100	0	0		
01111	0	1	0	z
10000	0	0		z
10001	0	1	0	z
10100	0	0		
10111	0	1	0	z
11000	1	1	$\bar{z}+\bar{z}$	$\bar{z}+z\bar{d}$
11011	x	x	$\bar{z}d$	$\bar{z}+z\bar{d}$
11100	x	x	d	d
11111	x	x	d	d
00000	x	x	d	d
00001	x	x		
00110	x	x	$\bar{z}d$	$\bar{z}d$
00111	0	0		
01000	1	1	$z+\bar{z}$	\bar{z}
01011	1	0	$z+\bar{z}$	\bar{z}
01110	1	0	$z+\bar{z}$	0
01111	1	0	$z+\bar{z}$	0
10000	1	0	$z+\bar{z}$	0
10001	1	0	$z+\bar{z}$	0
10110	1	0	$z+\bar{z}$	0
10111	1	0	$z+\bar{z}$	0
11000	1	0	$\bar{z}+\bar{z}$	$\bar{z}d$
11011	x	x	$\bar{z}d$	$z\bar{d}$
11110	x	x	d	d
11111	x	x	d	d

5.



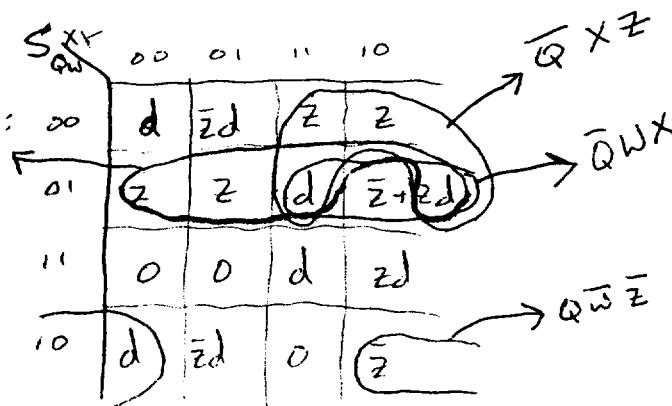
$$D = \bar{Q}\bar{W}\bar{X} + wx + QW + QX$$

$$S = \bar{Q}XZ + \bar{Q}WX + QWZ + Q\bar{W}\bar{Z}$$

Minimal SOP, solved separately

$$D = \bar{Q}\bar{W}\bar{X} + \bar{Q}WX + QW + QX$$

$$S = \text{same as above}$$

Minimal gate count, by using QWX term from S twice, i.e. reusing it in D .

6. Make a Store-Toggle flip flop with the following excitation table

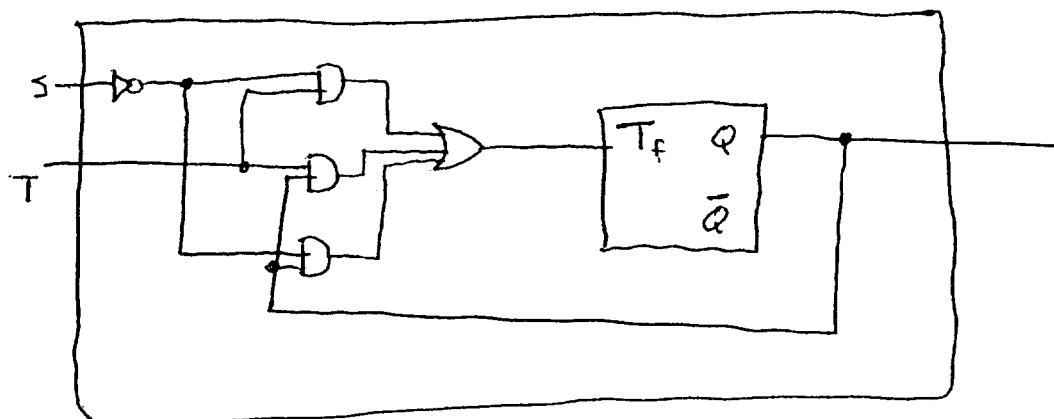
S	T	Q^{t+1}
0	0	0 reset
0	1	\bar{Q}^t Toggle
1	0	Q^t store
1	1	1 reset

$Q^t S T$	Q^{t+1}	T_f
000	0	0
001	1	1
010	0	0
011	0	0
100	0	1
101	0	1
110	1	0
111	0	1

T_f	S	00	01	11	10
0	0	0	1	0	0
1	1	1	1	1	0

$$T_f = \bar{S}T + Q^t \bar{S} + Q^t T$$

~~Don't care states are not considered~~



Q^t	Q^{t+1}	ST
0	0	Anything except 01
0	1	01
1	0	Anything except 10
1	1	10

2 bits of domain control

Transfer x if $\overline{S_A}$ count down

Call $S_A = A$ $S_B = B$ $\overline{S_A} \leftarrow R \oplus 1$

Initial
Program

$A^T B^T \times$	$A^T B^T$	$S_A \overline{R_A}$	$S_B \overline{R_B}$
0 0 0 0	1 1	1 0	1 0
0 0 0 1	0 1	0 x	0 0
0 0 1 0	0 0	0 x	0 1
0 0 1 1	0 1	0 0	0 1
0 1 0 0	0 1	1 0	1 0
0 1 0 1	0 0	1 x	1 0
0 1 1 0	0 1	0 0	1 1
0 1 1 1	0 0	1 x	1 1
1 0 0 0	0 1	0 0	0 0
1 0 0 1	0 0	0 x	0 0
1 0 1 0	0 1	0 0	0 1
1 0 1 1	0 0	0 x	0 1
1 1 0 0	0 1	0 0	0 0
1 1 0 1	0 0	0 x	0 0
1 1 1 0	0 1	0 0	0 0
1 1 1 1	0 0	0 x	0 0

$G^T A^T \leftarrow S_B$

$G^T A^T \leftarrow S_A$

$G^T A^T \leftarrow S_B$

$G^T A^T \leftarrow S_A$

S_A

$$\begin{array}{c} S_A \\ \diagup \quad \diagdown \\ A^T \quad B^T \\ \hline S_A = \frac{\overline{A^T B^T} \cdot \overline{x} + \overline{A^T B^T} \cdot x}{A^T \otimes B^T \otimes \overline{x}} \end{array}$$

$$S_B = A^T \overline{B^T} \cdot \overline{x} + A^T \overline{B^T} \cdot x$$

$$\begin{array}{c} \diagup \quad \diagdown \\ A^T \quad B^T \\ \hline S_B = A^T \otimes B^T \otimes x \end{array}$$

$$F_A = \overline{S_A}$$

$$\begin{array}{c} \diagup \quad \diagdown \\ A^T \quad B^T \\ \hline F_A = A^T \otimes B^T \otimes \overline{x} \end{array}$$

$$\begin{array}{c} \diagup \quad \diagdown \\ A^T \quad B^T \\ \hline F_B = B^T \end{array}$$

$$\begin{array}{c} \diagup \quad \diagdown \\ A^T \quad B^T \\ \hline F_A = \overline{S_A} \quad F_B = B^T \end{array}$$

$$\begin{array}{c} \diagup \quad \diagdown \\ A^T \quad B^T \\ \hline F_A = \overline{S_A} \quad F_B = B^T \end{array}$$

8. Design a 4-bit ring counter with initialize and error correction. The only inputs are the clock, and I , the initialize signal.
The system outputs are:-

1000
 0100
 0010
 0001
 1000
 :

If the system enters any invalid state, it must return to 1000 on the next clock cycle. The signal to do this is called E .

Use D flip flops : $D_3 D_2 D_1 D_0$.

Just write the equations for each D input - you don't need to draw them. Also write the equation for E .

$E = \text{You don't have an odd # of 1's} + \text{You have 3 1's}$

$$\therefore E = \overline{Q_4 \oplus Q_3 \oplus Q_2 \oplus Q_1} + Q_3 Q_2 Q_1 + Q_4 Q_2 Q_1 + Q_4 Q_3 Q_1 + Q_4 Q_3 Q_2$$

$$D_3 = I + E + Q_0$$

$$D_2 = \overline{I} \overline{E} Q_3$$

$$D_1 = \overline{I} \overline{E} Q_2$$

$$D_0 = \overline{I} \overline{E} Q_1$$

$$E = Q_3 Q_2 Q_1 Q_0 + Q_3 Q_2$$

$$9. \quad F = A \bar{B} \bar{C} \bar{D} + \bar{B} + C$$

